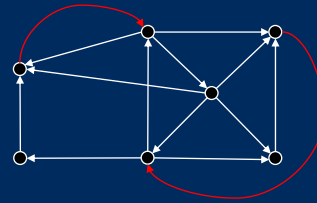


Fourier Analysis with Direction

M. Püschel



Department of Computer Science

ETH zürich

1

An Approach to Generalizing Signal Processing

M. Püschel and J. M. F. Moura
[Algebraic Signal Processing Theory](#)
<http://arxiv.org/abs/cs.IT/0612077>, 2006

M. Püschel and J. M. F. Moura
[Algebraic Signal Processing Theory: Foundation and 1-D Time](#)
IEEE TSP, 2008

M. Püschel and J. M. F. Moura
[Algebraic Signal Processing Theory: 1-D Space](#)
IEEE TSP, 2008

J. Kovacevic and M. Püschel
[Algebraic Signal Processing Theory: Sampling for Infinite and Finite 1-D Space](#)
IEEE TSP, 2010

M. Püschel and M. Rötteler
[Algebraic Signal Processing Theory: 2-D Spatial Hexagonal Lattice](#)
IEEE TIP, 2007

A. Sandryhaila, J. Kovacevic and M. Püschel
[Algebraic Signal Processing Theory: 1-D Nearest-Neighbor Models](#)
IEEE TSP, 2012

M. Püschel and M. Rötteler
[Fourier Transform for the Directed Quincunx Lattice](#)
Proc. ICASSP, 2005

M. Püschel and M. Rötteler
[Fourier Transform for the Spatial Quincunx Lattice](#)
Proc. ICASSP, 2005

Concise Derivation of Fast Transforms

M. Püschel and J. M. F. Moura
[Algebraic Signal Processing Theory: Cooley-Tukey Type Algorithms for DCTs](#)
IEEE TSP, 2008

Y. Voronenko and M. Püschel
[Algebraic Signal Processing Theory: Cooley-Tukey Type Algorithms for Real DFTs](#)
IEEE TSP, 2009

M. Püschel and M. Rötteler
[Algebraic Signal Processing Theory: Cooley-Tukey Type Algorithms on the 2-D Spatial Hexagonal Lattice](#)
Applicable Algebra, 2008

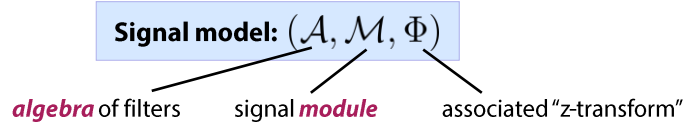
A. Sandryhaila, J. Kovacevic and M. Püschel
[Algebraic Signal Processing Theory: Cooley-Tukey Type Algorithms for Polynomial Transforms Based on Induction](#)
SIAM JMAA, 2011

2

2

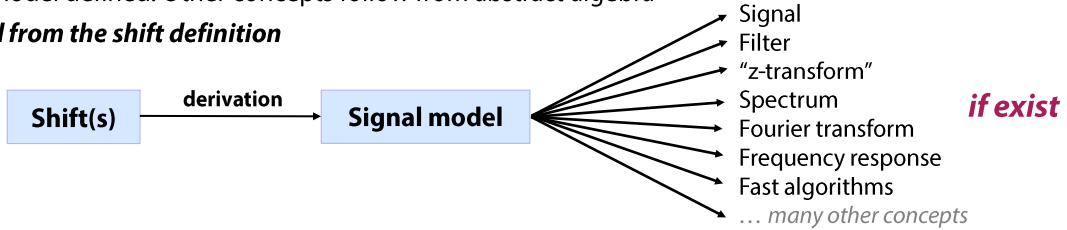
Key Insights

Axiomatic foundation



Signal model defined: Other concepts follow from abstract algebra

Derived from the shift definition

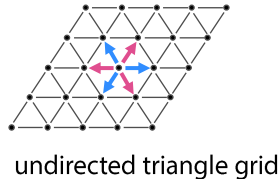
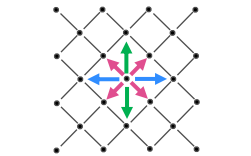
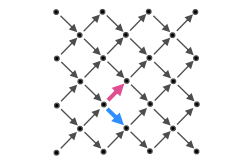
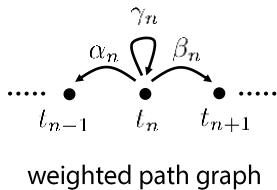
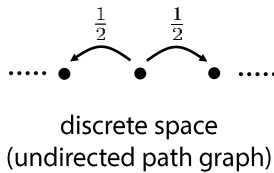
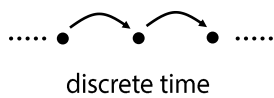


The shift is everything!

$(\mathcal{A}, \mathcal{M}, \Phi)$ shift invariant $\iff \mathcal{A}$ is commutative = a polynomial algebra

Detailed insight into boundary conditions: why needed and which are good

Considered Shifts Back Then (boundary conditions omitted)



M. Püschel and J. M. F. Moura
Algebraic Signal Processing Theory
<http://arxiv.org/abs/cs.IT/0612077>, 2006

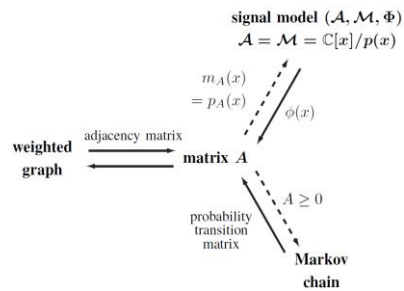


Fig. 22. The connection between square matrices A , weighted graphs, regular signal models, and finite Markov chains.

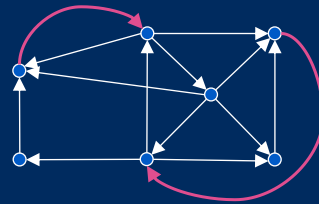
Every weighted graph defines a signal model
Signal model with one shift = (adjacency) GSP

*“Detailed insight into boundary conditions:
why needed and which are good”*

Part 1: Fourier Analysis on Directed Graphs



with
Bastian Seifert



Digraph Signal Processing with Generalized Boundary Conditions, IEEE TSP 2021

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Graph Signal Processing...

Concept	Undirected Graphs
Shift/Variation operator	✓ (Adjacency or Laplacian) Symmetric
Convolution	✓
Fourier Basis/Transform (eigendecomposition)	✓
Orthogonality	✓

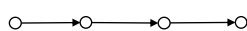
6

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Graph Signal Processing...

Concept	Undirected Graphs	Directed Graphs (Digraphs)
Shift/Variation operator	✓ (Adjacency or Laplacian) Symmetric	✓ (Adjacency or Laplacian) Not symmetric
Convolution	✓	✓
Fourier Basis/Transform (eigendecomposition)	✓	✗ Does not exist in general
Orthogonality	✓	✗ (In general no)

Digraph Example:



$$A = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}, \quad L = D - A = \begin{bmatrix} 1 & -1 & 0 & 0 \\ 0 & 1 & -1 & 0 \\ 0 & 0 & 1 & -1 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

No Eigendecomposition/Fourier basis

Only one eigenvalue (Jordan Block)

7

7

Digraph Signal Processing is an Open Problem

Graph Signal Processing: Overview, Challenges, and Applications, Ortega et al., Proc. IEEE, vol. 106(5), pp. 808–828, 2018

Jordan normal form + tricks:

Spectral Projector-Based Graph Fourier Transforms

Joya A. Deri, Member, IEEE, and José M. F. Moura, Fellow, IEEE

Agile Inexact Methods for Spectral Projector-Based Graph Fourier Transforms

Joya A. Deri, Member, IEEE, and José M. F. Moura, Fellow, IEEE

Graph Fourier Transform based on Directed Laplacian

Rahul Singh, Abhishek Chakraborty, Graduate Student Member, IEEE, and B. S. Manoj, Senior Member, IEEE

Fourier transform via optimization:

On the Graph Fourier Transform for Directed Graphs

Stefania Sardellitti, Member, IEEE, Sergio Barbarossa, Fellow, IEEE, and Paolo Di Lorenzo, Member, IEEE

A Directed Graph Fourier Transform With Spread Frequency Components

Rasoul Shefipour¹, Student Member, IEEE, Ali Khodabakhsh², Student Member, IEEE, Gonzalo Mateos³, Senior Member, IEEE, and Evdokia Nikolova

DIGRAPH FOURIER TRANSFORM VIA SPECTRAL DISPERSION MINIMIZATION

Rasoul Shefipour¹, Ali Khodabakhsh², Gonzalo Mateos³, and Evdokia Nikolova²

Approximate diagonalization:

Graph Fourier Transform: A Stable Approximation

Jolo Domingos and José M. F. Moura

Complex-valued Shift:

Graph Signal Processing for Directed Graphs based on the Hermitian Laplacian

Satoshi Furutani¹, Toshiki Shilohara¹, Mitsunaki Akiyama¹, Kunio Hato³, and Masaki Aida²

ORTHOGONAL TRANSFORMS FOR SIGNALS ON DIRECTED GRAPHS

Julia Barrufet¹, Antonio Ortega²

Subset of diagonalizable filters:

DIAGONALIZABLE SHIFT AND FILTERS FOR DIRECTED GRAPHS BASED ON THE JORDAN-CHEVALLEY DECOMPOSITION

Panagiotis Misiakos*

Chris Wendler, Markus Püschel

8

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Our Key Idea: Boundary Conditions

Classical Signal Processing

Digraph for finite discrete time

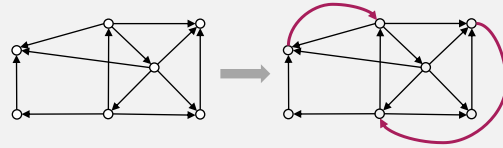


$$A + B = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$\text{DFT}_4(A + B) \text{DFT}_4^{-1} = \begin{bmatrix} 1 & & & \\ & i & & \\ & & -1 & \\ & & & -i \end{bmatrix}$$

This is done even though
the signal is usually not periodic!

Can we do this for other digraphs?



$$\begin{bmatrix} 0 & 1 & & & & & \\ 0 & 0 & & & & & \\ & 0 & 1 & & & & \\ & & 0 & 0 & & & \\ & & & 1 & & & \\ & & & & \omega_3^2 & & \\ & & & & & \omega_3 & \\ & & & & & & 1 \end{bmatrix} \rightarrow \begin{bmatrix} \lambda_1 & & & & & & \\ & \lambda_2 & & & & & \\ & & \lambda_3 & & & & \\ & & & \lambda_4 & & & \\ & & & & \lambda_5 & & \\ & & & & & \lambda_6 & \\ & & & & & & \lambda_7 \end{bmatrix}$$

Goal: Add small number of edges to obtain Fourier basis of eigenvectors

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Edges to Destroy Jordan Blocks

Tool: Perturbation Theory

On the Change in the Spectral Properties of a Matrix under Perturbations of Sufficiently Low Rank

S. V. Savchenko

Theorem 1. Let A be an arbitrary square matrix, and let $B = \sum_{i=1}^r (\cdot, \xi_i) \eta_i$ be an operator of rank r . Consider any eigenvalue λ of A . We arrange the sizes $n_1 \geq \dots \geq n_k$ of the corresponding Jordan blocks in nonascending order. Suppose that $k \geq r$ and

$$n_{A+B}(\lambda) = n_A(\lambda) - n_1 - \dots - n_r. \quad (4)$$

Then n_{r+1}, \dots, n_k are the sizes of Jordan blocks of the matrix $A+B$ associated with λ .

LOW RANK PERTURBATION OF JORDAN STRUCTURE*

JULIO MORO[†] AND FROILÁN M. DOPICO[†]
CONCLUDING THEOREM. Let A be a complex $n \times n$ matrix and λ_0 an eigenvalue of A with geometric multiplicity g . Let B be a complex $n \times n$ matrix with $\text{rank}(B) \leq g$ and C_0 be as in the statement of Theorem 2.1. Then the Jordan blocks of $A+B$ with eigenvalue λ_0 are just the $g - \text{rank}(B)$ smallest Jordan blocks of A with eigenvalue λ_0 if and only if $C_0 \neq 0$.

Our work: Specialize to Adjacency/Laplacian matrices

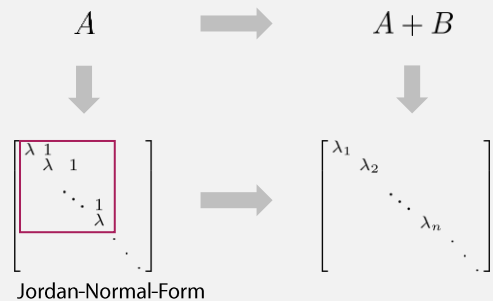
Corollary: Adding an edge is enough to **destroy** the largest **Jordan block** to a choosen eigenvalue.

Q: How to find these edges?

Theorem: Let $u_1, \dots, u_r, v_1, \dots, v_r$ be left/right eigenvectors of Jordan blocks to the eigenvalue λ and B the matrix containing only the new edge, then if

$$\sum_{k=1}^r u_k^T B v_k \neq 0$$

the largest Jordan block of λ gets destroyed in $A+B$.



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The Algorithm

```

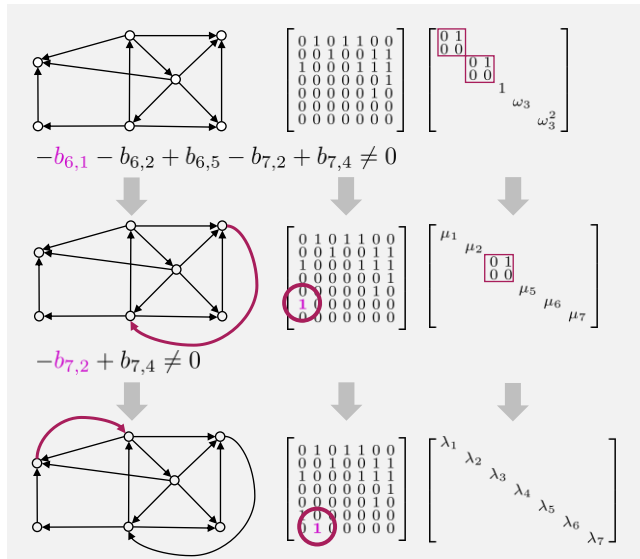
function DESTROYALLJORDANBLOCKS(A)
while A not diagonalizable do
  u1, ..., ur ← left EVs to largest Jordan blocks
  v1, ..., vr ← right EVs to largest Jordan blocks
  if ∃(i, j) s.t. ∑k uk,jvk,i ≠ 0 and Ai,j = 0 then
    Ai,j ← 1
  else
    select (i, j) random s.t. Ai,j = 0
    Ai,j ← 1
  end if
end while
return A
end function
    
```

Numerical Implementation:
Needs some tricks, see paper

Faster, Inexact Implementation:
Eigenvalue 0 appears often, use that

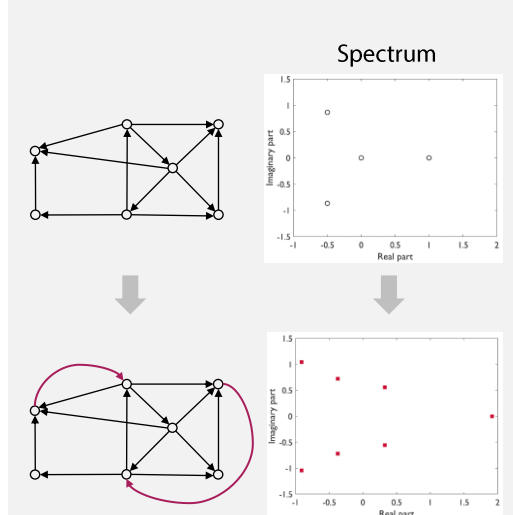
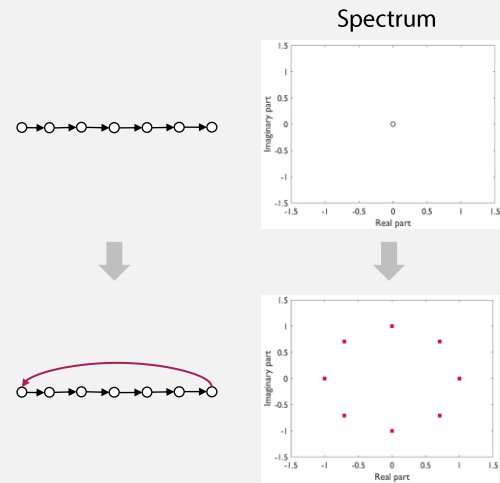
Variants: Destroy zero eigenvalues so
Fourier transform is numerically stable

Also works with Laplacian D – A ...



Behaviour of Eigenvalues? (Spectrum)

Finite Discrete Time



Generally Applicable & Fast

Random digraphs with different properties, **500 nodes & ~5000 edges**

	min		median		max	
	edges	time	edges	time	edges	time
Watts-Strogatz	0	0.2s	1	0.5s	3	1.3s
Barabási-Albert	36	4.4s	44	10s	55	31s
Klemm-Eguílez	10	2.2s	27	6s	47	9s

Medium number of edges added: 27
Median runtime: 6 seconds

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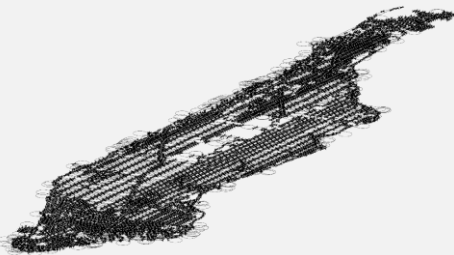
13

Scalable

Manhattan Taxi Graph

Li & Moura, ECAI, 2020

5464 nodes & 11568 edges

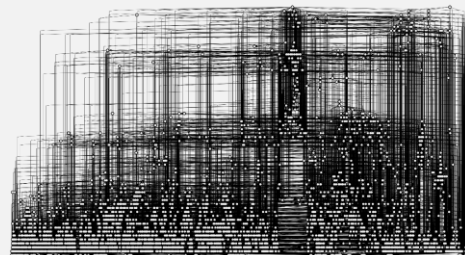


Runtime: 19 hours,
243 edges added
Runtime (inexact algo): **5 min**,
772 edges added

Citation Graph

<https://snap.stanford.edu/data/cit-HepPh.html>

4989 nodes & 17840 edges
(almost acyclic)



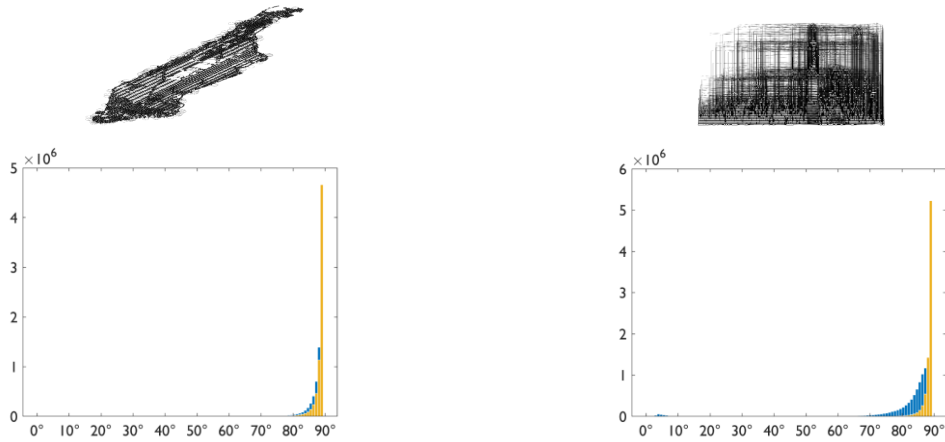
Runtime (inexact algo): **31.5 min**,
1911 edges added

14

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Fourier Bases Found are Almost Orthogonal

Histograms of pairwise angles between basis vectors



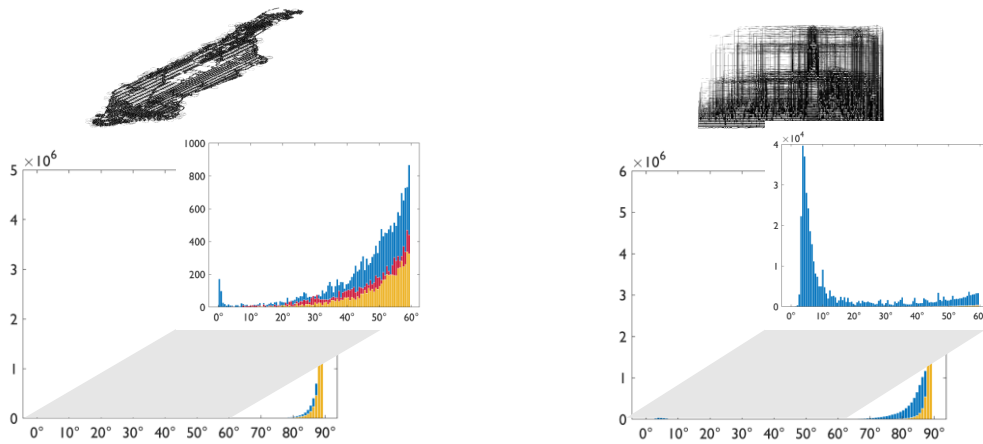
and numerically stable (needed for inversion to obtain Fourier transform)

15

15

Fourier Bases Found are Almost Orthogonal

Histograms of pairwise angles between basis vectors



and numerically stable (needed for inversion to obtain Fourier transform)

16

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One Application: Denoising with Neural Networks

S. Chen, Y. C. Eldar, and L. Zhao, **Graph Unrolling Networks: Interpretable Neural Networks for Graph Signal Denoising**, IEEE TSP 2021



ported to digraphs (needs a Fourier basis)

V. Mihal, B. Seifert, M. Püschel, **Porting Signal Processing from Undirected to Directed Graphs: Case Study Signal Denoising with Unrolling Networks**, Proc. EUSIPCO 2022

	ForgetDir	AdHoc	SelfLoops	UndShift	AssocShift
$p = 0.5$					
SprFreq	0.43 ± 0.01	0.46 ± 0.01	0.35 ± 0.01	0.19 ± 0.00	-
HermL	0.28 ± 0.01	0.31 ± 0.01	0.19 ± 0.00	0.07 ± 0.01	0.19 ± 0.04
StabApp	0.13 ± 0.00	0.08 ± 0.00	0.09 ± 0.00	0.10 ± 0.00	-
GenBC1	0.10 ± 0.00	0.01 ± 0.00	0.01 ± 0.00	0.03 ± 0.00	0.01 ± 0.00
GenBC2	0.10 ± 0.00	0.01 ± 0.00	0.01 ± 0.00	0.03 ± 0.00	0.01 ± 0.00
$p = 0.9$					
SprFreq	0.19 ± 0.01	0.19 ± 0.01	0.12 ± 0.00	0.06 ± 0.00	-
HermL	0.21 ± 0.01	0.28 ± 0.01	0.16 ± 0.01	0.07 ± 0.00	0.19 ± 0.01
StabApp	0.24 ± 0.00	0.10 ± 0.00	0.11 ± 0.00	0.11 ± 0.00	-
GenBC1	0.15 ± 0.01	0.03 ± 0.00	0.03 ± 0.01	0.06 ± 0.01	0.05 ± 0.00
GenBC2	0.15 ± 0.00	0.06 ± 0.00	0.06 ± 0.00	0.05 ± 0.00	0.01 ± 0.00

works well

TABLE I: NMSE results of denoising low-frequency signals.

17

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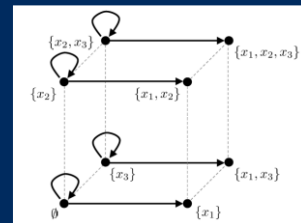
"shift defines signal model"

"SP with one shift = GSP"

Part 2.1: Fourier Analysis on Powersets



with
Chris Wendler



Discrete Signal Processing with Set Functions, IEEE SP 2021
Learning Set Functions that are Sparse in Non-Orthogonal Fourier Bases, Proc. AAAI 2021

18

Fourier Meets Möbius: Fast Subset Convolution

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Proc. STOC 2007

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ABSTRACT

We present a fast algorithm for the *subset convolution problem*: given functions f and g defined on the lattice of subsets of an n -element set N , compute their *subset convolution* $f * g$, defined for all $S \subseteq N$ by

$$(f * g)(S) = \sum_{T \subseteq S} f(T)g(S \setminus T),$$

shift?
Fourier transform?
frequency ordering?
frequency response?
applications?

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Signals on Powersets = Set Functions

Finite set: $N = \{x_1, \dots, x_n\}$

Its power set: $2^N = \{A \mid A \subseteq N\}$

Set function: $s = (s_A)_{A \subseteq N}$

Shift by $x_i \in N$ $s \mapsto (s_{A \setminus \{x_i\}})$

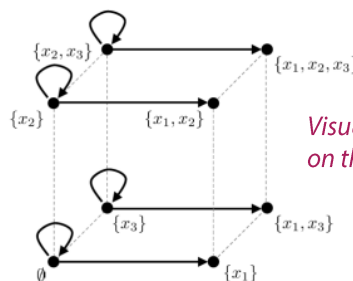
"Delay by x_i "

not invertible

n shifts!

A. Krause and D. Golovin, **Tractability: Practical Approaches to Hard Problems**. ch. **Submodular function maximization**, pp. 71–104, 2014

Applications: Recommender systems, image processing, facility location, online auctions etc.



Visualization of shift by x_1 on the associated "z-transform"

Signal domain = directed hypercube

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Derived Signal Model and SP Concepts

Shift(s) by $x_i \in N$ $s \mapsto (s_{A \setminus \{x_i\}})$ *n shifts commute: shift-invariance*

Convolution (filter): $(h * s)_A = \sum_{X \subseteq N} h_X s_{X \setminus A}$

Fourier basis (as matrix): $\begin{bmatrix} 1 & 0 \\ 1 & -1 \end{bmatrix} \otimes \dots \otimes \begin{bmatrix} 1 & 0 \\ 1 & -1 \end{bmatrix}$

Fourier transform: $\hat{s}_B = \sum_{A \subseteq B} (-1)^{|A|} s_A$

as matrix: DSFT $_{2^n} = \begin{bmatrix} 1 & 0 \\ 1 & -1 \end{bmatrix} \otimes \dots \otimes \begin{bmatrix} 1 & 0 \\ 1 & -1 \end{bmatrix}$ *diagonalizes all shifts and filters
not orthogonal, lower triangular
 $n2^{n-1}$ ops*

Frequency response: $\bar{h}_B = \sum_{A \subseteq N, A \cap B = \emptyset} (-1)^{|A|} h_A$

Convolution theorem: $\widehat{h * s} = \bar{h} \odot \hat{s}$

21

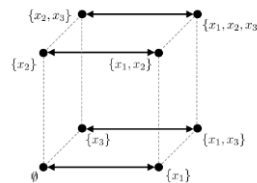
21

Four Variants of Signal Models

model	q, A	on signal	Q, A	on signal	$(h * s)_A$	matrix for q
1	$A \cup \{q\}$	$s_A + s_{A \setminus \{q\}}, q \in A$ 0, else	$A \cup Q$	$\sum_{A \setminus Q \subseteq B \subseteq A} s_B, Q \subseteq A$ 0, else	$\sum_{Q \cup B = A} h_Q s_B$	$\begin{bmatrix} 0 & 0 \\ 1 & 1 \end{bmatrix}$
2	$A \setminus \{q\}$	$s_A + s_{A \cup \{q\}}, q \notin A$ 0, else	$A \setminus Q$	$\sum_{B \subseteq Q} s_{A \cup B}, Q \subseteq N \setminus A$ 0, else	$\sum_{Q \subseteq N \setminus A} \sum_{B \subseteq Q} h_Q s_{A \cup B}$	$\begin{bmatrix} 1 & 1 \\ 0 & 0 \end{bmatrix}$
3	$A + A \cup \{q\}, q \notin A$ 0, else	$s_{A \setminus \{q\}}$	$\sum_{B \subseteq Q} A \cup B, Q \subseteq N \setminus A$ 0, else	$s_{A \setminus Q}$	$\sum_{Q \subseteq N} h_Q s_{A \setminus Q}$	$\begin{bmatrix} 1 & 0 \\ 1 & 0 \end{bmatrix}$
4	$A + A \setminus \{q\}, q \in A$ 0, else	$s_{A \cup \{q\}}$	$\sum_{B \subseteq Q} A \setminus B, Q \subseteq A$ 0, else	$s_{A \cup Q}$	$\sum_{Q \subseteq N} h_Q s_{A \cup Q}$	$\begin{bmatrix} 0 & 1 \\ 0 & 1 \end{bmatrix}$
5	$A \setminus \{q\} \cup \{q\} \setminus A =$ $A \cup \{q\}, q \notin A$ $A \setminus \{q\}, q \in A$	$s_{A \setminus \{q\} \cup \{q\} \setminus A} =$ $s_{A \cup \{q\}}, q \notin A$ $s_{A \setminus \{q\}}, q \in A$	$A \setminus Q \cup Q \setminus A$	$s_{A \setminus Q \cup Q \setminus A}$	$\sum_{Q \subseteq N} h_Q s_{A \setminus Q \cup Q \setminus A}$	$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$

Classical model for set functions
Fourier transform = Walsh-Hadamard transform

shift(s): $s \mapsto (s_{A \setminus \{x_i\}} \cup s_{\{x_i\} \setminus A})_{A \subseteq N}$



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Meaning of Spectrum?

Values of sets of goods

$$N = \text{[Laptop]} \text{ [Phone]} \text{ [Pen]}$$

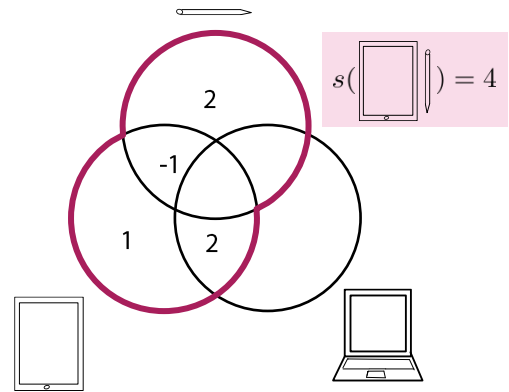
$$s(\text{[Phone]}) = 2 \quad s(\text{[Pen]}) = 1 \quad s(\text{[Laptop]}) = 2$$

$$s(\text{[Phone] [Pen]}) = 4 \quad \text{complementarity } 4 > 3$$

$$s(\text{[Phone] [Laptop]}) = 2 \quad \text{substitutability } 2 < 4$$

etc.

Weighted Venn Diagram of Values



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Fourier Transform



Values of all subsets

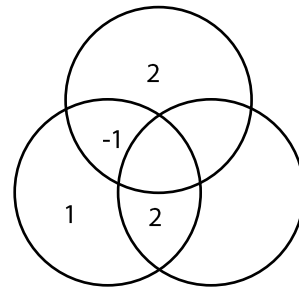


Values of fragments (interactions)

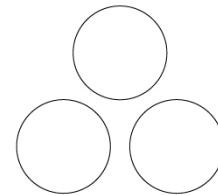
low frequencies

high frequencies

Spectrum is also a set function!



Fourier sparse



Very low frequency
(values of subsets = sums of values of items)

$$\hat{s}_B = 0, \quad |B| > 1$$

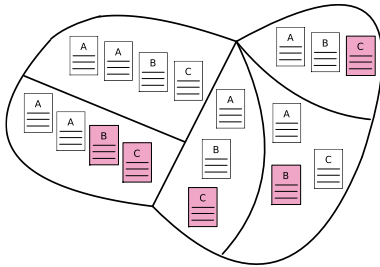
24

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Application: Preference Elicitation in Spectrum Auctions

N = licenses of bands electromagnetic spectrum

Set functions = bidders



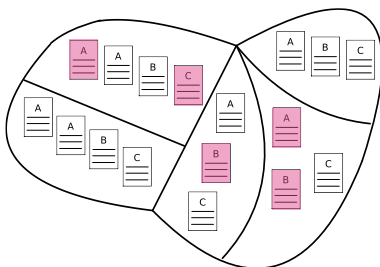
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Application: Preference Elicitation in Spectrum Auctions

N = licenses of bands electromagnetic spectrum

Set functions = bidders



Preference elicitation: estimate bidders from few (500 say) samples (values of subsets)

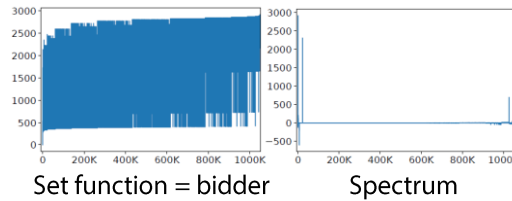
Idea: assume Fourier sparsity

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Set Function Estimate with Fourier Sparsity

Model bidder $|N| = 20$ licenses:



C. Wendler, A. Amrollahi, B. Seifert, A. Krause, M. Püschel
[Learning Set Functions that are Sparse in Non-Orthogonal Fourier Bases](#)
 Proc. AAAI, 2021

B. type	number of queries (in thousands)			Fourier coefficients recovered			relative reconstruction error		
	SSFT	SSFT+	H-WHT	SSFT	SSFT+	H-WHT	SSFT	SSFT+	H-WHT
local	3 ± 4	229 ± 73	781 ± 0	118 ± 140	303 ± 93	675 ± 189	0.5657 ± 0.4900	0 ± 0	0 ± 0
regional	20 ± 1	646 ± 12	781 ± 0	659 ± 32	813 ± 36	1,779 ± 0	0.0118 ± 0.0071	0 ± 0	0 ± 0
national	71 ± 0	3,305 ± 1	781 ± 0	1,028 ± 3	1,027 ± 6	4,170 ± 136	0.0123 ± 0.0014	0.0149 ± 0.0089	0.2681 ± 0.2116

$|N| = 98$

good reconstruction

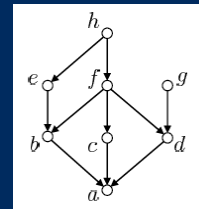
Later: collaboration with auction experts J. Weissteiner, C. Wendler, S. Seuken, B. Lubin, M. Püschel
[Fourier Analysis-based Iterative Combinatorial Auctions](#)
 Proc. International Joint Conference on Artificial Intelligence (IJCAI), 2022

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“shift defines signal model”

“SP with one shift = GSP”



Part 2.2: Fourier Analysis: From Powersets to Lattices



with
 Bastian Seifert
 Chris Wendler
 Tommaso Pegolotti

Discrete Signal Processing on Meet/Join Lattices, IEEE TSP 2021
Fast Möbius and Zeta Transforms, Arxiv

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Generalization to (Meet) Lattices

Power set convolution:

$$\begin{aligned}
 (\mathbf{h} * \mathbf{s})_A &= \sum_{X \subseteq N} h_X s_{X \setminus A} \\
 &= \sum_{X \subseteq N} h_X s_{X \cap N \setminus A} \\
 &\approx \sum_{X \subseteq N} h_X s_{X \cap A}
 \end{aligned}$$



Power set is an example of a meet lattice
meet operation = \cap

Meet lattice convolution:

$$(\mathbf{h} * \mathbf{s})_a = \sum_{x \in \mathcal{L}} h_x s_{x \wedge a}$$



639 pages
No convolution or
Fourier transform

Prior work:

A. Björklund, T. Husfeldt, P. Kaski, M. Koivisto, J. Nederlof, and P. Parviainen, **Fast zeta transforms for lattices with few irreducibles**, ACM Trans. on Algorithms, 2015.

(Meet) Lattices

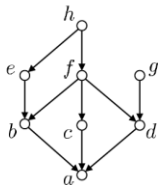
A lattice \mathcal{L} is a (here finite) poset: set with a partial order \leq

For $a, b, c \in \mathcal{L}$:

1. $a \leq a$,
2. $a \leq b$ and $b \leq a$ implies $a = b$, and
3. $a \leq b$ and $b \leq c$ implies $a \leq c$

For each $a, b \in \mathcal{L}$ there is a unique greatest lower bound $a \wedge b$

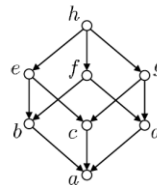
1. $a \wedge a = a$,
2. $a \wedge b = b \wedge a$, and
3. $(a \wedge b) \wedge c = a \wedge (b \wedge c)$



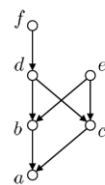
Some lattice



Total order lattice



Subset lattice (hypercube)



Not a lattice

Lattices \subset DAGs (directed acyclic graphs)

Derived Signal Model and SP Concepts

Lattice signal: $\mathbf{s} = (s_x)_{x \in \mathcal{L}}$

Shift(s) by $a \in \mathcal{L}$ $\mathbf{s} \mapsto (s_{x \wedge a})_{x \in \mathcal{L}}$

*"x gets delayed by a"
shifts commute: shift-invariance*

Convolution (filter): $(\mathbf{h} * \mathbf{s})_a = \sum_{x \in \mathcal{L}} h_x s_{x \wedge a}$

Inverse Fourier transform: $s_x = \sum_{y \leq x} \hat{s}_y$

*Fourier basis consists of 0/1 vectors
IFT sums spectral values of all predecessors*

Fourier transform: $\hat{s}_y = \sum_{x \leq y} \mu(x, y) s_x$

*diagonalizes all shifts and filters
not orthogonal, lower triangular*

$\mu(x, x) = 1$, *Möbius function*
 $\mu(x, y) = - \sum_{x \leq z < y} \mu(x, z)$, $x \neq y$

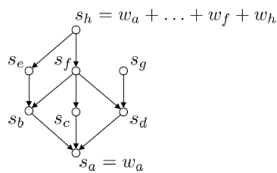
Spectrum: is partially ordered isomorphic to \mathcal{L}

G.-C. Rota, **On the foundations of combinatorial theory. I. theory of Möbius functions**,
Z. Wahrscheinlichkeitstheorie und Verwandte Gebiete, 1964.

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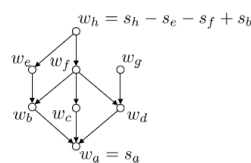
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Example



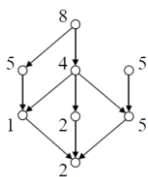
Signal

$s_h = w_a + w_b + w_c + w_d + w_e + w_f + w_h$
 $s_g = w_a + w_d + w_g$
 $s_f = w_a + w_b + w_c + w_d + w_f$
 $s_e = w_a + w_b + w_e$
 $s_d = w_a + w_d$
 $s_c = w_a + w_c$
 $s_b = w_a + w_b$
 $s_a = w_a$

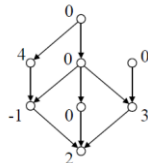


Spectrum

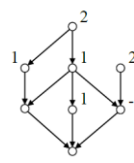
$w_h = s_h - s_f - s_e + s_b$
 $w_g = s_g - s_d$
 $w_f = s_f - s_d - s_c - s_b + 2s_a$
 $w_e = s_e - s_b$
 $w_d = s_d - s_a$
 $w_c = s_c - s_a$
 $w_b = s_b - s_a$
 $w_a = s_a$



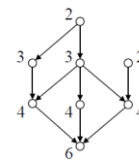
signal



spectrum



basic low pass filter
1 + sum of all shifts



its frequency response

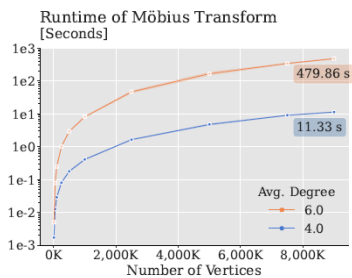
Applications: see paper

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Fast Lattice Fourier Transform

Can be computed in $O(nk)$, k = longest antichain in lattice

Enables computation for millions of nodes if DAG is sparse = low average degree



+ close to linear speedup when parallelized

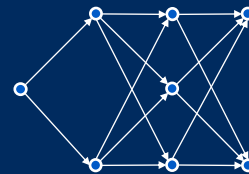
T. Pegolotti, B. Seifert, M. Püschel, [Fast Möbius and Zeta Transforms](#), Arxiv

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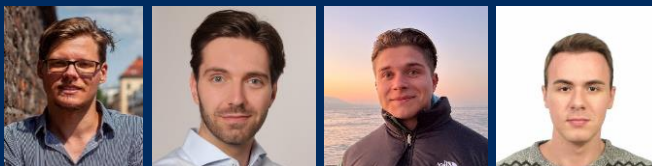
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"shift defines signal model"

"SP with one shift = GSP"



Part 2.3: Fourier Analysis: From Lattices to DAGs



with
Bastian Seifert
Chris Wendler
Panos Misiakos
Vedran Mihal

Causal Fourier Analysis on Directed Acyclic Graphs and Posets, Arxiv
Learning DAGs from Data with Few Root Causes, Arxiv

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Lattices once more:

Convolution $(\mathbf{h} * \mathbf{s})_a = \sum_{x \in \mathcal{L}} h_x s_{x \wedge a}$

Inverse Fourier transform: $s_x = \sum_{y \leq x} \hat{s}_y$

Fourier transform: $\hat{s}_y = \sum_{x \leq y} \mu(x, y) s_x$



Chris Wendler

“Uses only the partial order. I think we don’t need the lattice property. Partial order/DAG is enough”

But how to do the shift without \wedge ?

Also, we need weighted edges for broader applicability

DAGs are, in a sense, “worst case” in GSP:

But highly relevant in causal reasoning, Bayesian networks etc.

$$\begin{bmatrix} 0 & & & & \\ & \ddots & & & \\ & & \ddots & & \\ & & & \ddots & \\ 0 & \dots & & & 0 \end{bmatrix} \text{ adjacency matrix only } 0 \text{ eigenvalue}$$

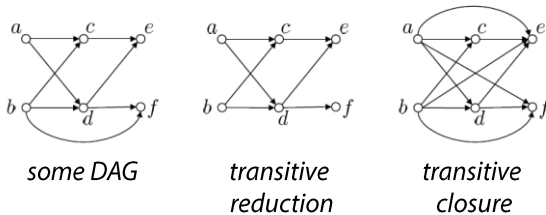
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DAGs \approx Posets

Every DAG defines a unique partial order: $a \leq b \iff a$ is a predecessor of b in the DAG

A partial order can be represented by several DAGs:



all define the same partial order

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Signal Model for DAGs

Weighted DAG: $D = (V, E, A)$ with induced \leq

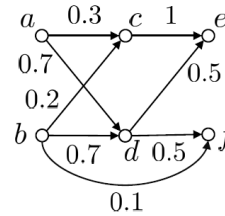
Observed signal on DAG: $s = (s_x)_{x \in V}$

The signal satisfies:
$$s_x = \sum_{y \leq x} w_{y,x} c_y$$

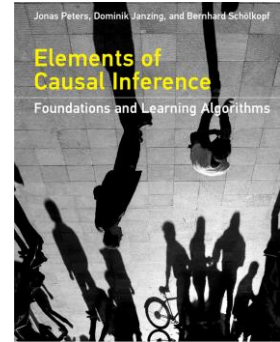
↑ Unknown value produced by each node that we call **cause**

↑ Weights
How obtained from A?

↑ Weighted version of the previous inverse Fourier transform on lattices



This equation does not imply causality

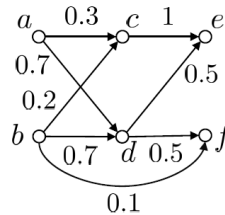


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$$s_x = \sum_{y \leq x} w_{y,x} c_y$$

$$s = Wc$$

↑ Weights
How obtained from A?



River network

Nodes: cities

Edges: rivers

Edge weights: fraction of pollution propagated

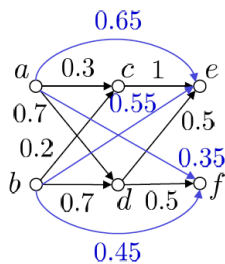
c_y unknown pollution inserted by city y

s_x measured pollution at city x

accumulated from all predecessor nodes

In this case: $W = I + A + A^2 + \dots + A^{n-1} = (I - A)^{-1}$

$(+, \cdot)$ -transitive closure of A



Fourier transform = $W^{-1} = I - A$:

$$\begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ -0.3 & -0.2 & 1 & 0 & 0 & 0 \\ -0.7 & -0.7 & 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & -0.5 & 1 & 0 \\ 0 & -0.1 & 0 & -0.5 & 0 & 1 \end{bmatrix}$$

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Transitive closures

S	$u \oplus v$	$u \odot v$	0_S	1_S	Meaning of edge weight $\bar{a}_{x,y}$ in closure
$\{0, 1\}$	u or v	u and v	0	1	Reachability in graphs; x is cause of y
$[0, 1]$	$u + v$	$u \cdot v$	0	1	Fraction of pollution from x reaching y ← prior slide
$[0, 1]$	$\max(u, v)$	$u \cdot v$	0	1	Strongest influence/most reliable path from x to y
$\mathbb{R}^+ \cup \{\infty\}$	$\min(u, v)$	$u + v$	∞	0	Shortest path length from x to y
$\mathbb{R}^+ \cup \{\infty\}$	$\max(u, v)$	$\min(u, v)$	0	∞	Largest capacity path from x to y

S. K. Abdali and B. D. Saunders, **Transitive closure and related semiring properties via eliminants**, Theor. Comput. Sci., 1985

$O(n^3)$ computation with generic Floyd-Warshall algorithm

```

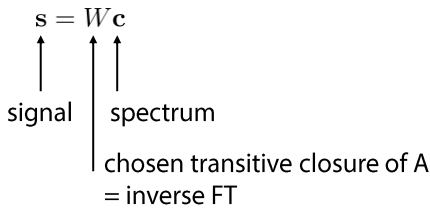
function WEIGHTEDTRANSITIVECLOSURE(W)
    H(0) ← A
    for k = 1, ..., n do
        for i = 1, ..., n do
            for j = 1, ..., n do
                h(k)(xi, xj) ← h(k-1)(xi, xj) ⊕ (h(k-1)(xi, xk) ⊙ h(k-1)(xk, xj))
            end for
        end for
    end for
    return Ā = H(n)
end function
    
```

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What About the Shift?

$$s_x = \sum_{y \leq x} w_{y,x} c_y$$

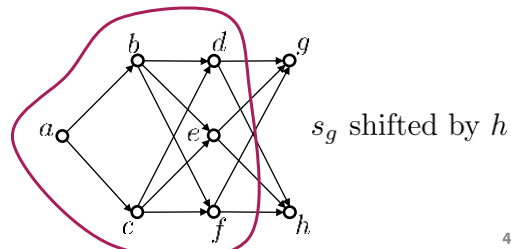
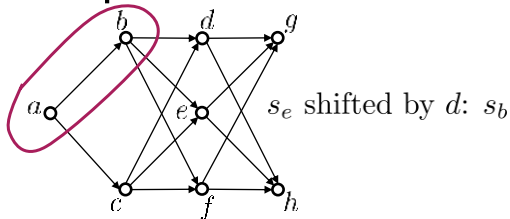


Shift(s) by $q \in V$

$$s_x = \sum_{y \leq x} w_{y,x} c_y \mapsto \sum_{y \leq x, y \leq q} w_{y,x} c_y$$

In frequency domain:
 Sum over common causes of x and q
 =
 Multiplies non-common causes by 0
 Others by 1

Example:



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Experiments

Fourier sparsity
=
Few causes

*e.g., few cities pollute
in one data set*

Learn DAG signal from samples under the assumption of Fourier sparsity
Example infection data: see paper

Learn the DAG from data under the assumption of Fourier sparsity: next

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Linear Structural Causal Models

Linear SCM (also SEM) with DAG A : assumes data generation as

$$\mathbf{s} = A\mathbf{s} + \mathbf{n}$$

\uparrow iid Gaussian noise

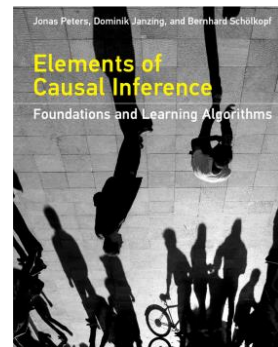
*can represent any n -variate
Gaussian distribution*

In our language:

$$\iff \mathbf{s} = W\mathbf{n}, \quad W = (I - A)^{-1}$$

\uparrow dense, random spectrum

$(+, \cdot)$ transitive closure of A



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Learning DAG from Data

Classical linear SCM model: $\mathbf{s} = W\mathbf{n}, \quad W = (I - A)^{-1}$

Our assumption (Fourier sparsity): $\mathbf{s} = W(\underbrace{\mathbf{c} + \mathbf{n}_c}_{\text{sparse plus noise}}) + \mathbf{n}_s$
measurement noise

Optimization problem to find DAG: given data matrix S

$$\min_{\mathbf{A} \in \mathbb{R}^{n \times n}} \|\mathbf{S}W^{-1}\|_1 + \lambda \|\mathbf{A}\|_1 \quad \text{s.t.} \quad \text{trace}(e^{\mathbf{A} \odot \mathbf{A}} - n) = 0$$

\uparrow
sparse spectrum
 \uparrow
sparse A
 \uparrow
A acyclic
(idea from NoTears)

X. Zheng, B. Aragam, P. K. Ravikumar, and E. P. Xing. **Dags with no tears: Continuous optimization for structure learning.** Advances in Neural Information Processing Systems, 2018

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Results

40 nodes, 80 edges
30% Fourier sparsity

Metric:

SHD = number of edge deletions/insertions/reversals to obtain true DAG

Hyperparameter	Default	Change	Varsort.	SparseRC (ours)	DAGMA	GOLEM	NOTEARS	DAG-NoCurl	GES	sortnregress
1. Default settings			0.96	0.00 ± 0.00	25.80 ± 3.71	2.70 ± 2.24	3.10 ± 2.30	28.20 ± 5.67	56.90 ± 25.84	failure
2. Graph type	Erdős-Renyi	Scale-free	0.99	0.00 ± 0.00	26.20 ± 4.79	1.10 ± 1.38	0.90 ± 0.94	7.60 ± 5.68	78.00 ± 39.52	33.30 ± 9.23
3. Edges / vertices	2	3	0.97	0.90 ± 1.04	44.40 ± 7.10	4.30 ± 2.87	10.40 ± 8.26	40.50 ± 14.02	failure	failure
4. Larger weights in A	(0.4, 0.8)	(0.5, 2)	0.98	23.10 ± 13.66	13.80 ± 3.63	1.60 ± 3.26	8.50 ± 8.42	20.30 ± 8.53	64.30 ± 35.96	failure
5. High sparsity in C	p = 0.3	p = 0.1	0.96	0.00 ± 0.00	65.70 ± 3.23	4.90 ± 4.28	9.80 ± 2.31	36.90 ± 7.78	47.40 ± 25.42	failure
6. Low sparsity in C	p = 0.3	p = 0.6	0.96	64.80 ± 4.40	14.30 ± 4.22	2.50 ± 1.96	2.50 ± 1.36	32.30 ± 11.44	63.60 ± 25.05	failure
7. N _c , N _s deviation	σ = 0.01	σ = 0.1	0.96	3.50 ± 1.75	21.00 ± 2.53	2.50 ± 1.36	3.50 ± 2.50	26.40 ± 8.91	50.80 ± 13.37	75.40 ± 19.27
8. N _c , N _s distribution	Gaussian	Gumbel	0.96	0.20 ± 0.40	27.00 ± 3.29	4.10 ± 2.98	4.60 ± 3.38	24.30 ± 8.06	53.10 ± 21.19	failure
9. Measurement noise	N _c ≠ 0, N _s = 0	N _c = 0, N _s ≠ 0	0.96	0.10 ± 0.30	24.10 ± 3.70	3.40 ± 4.15	2.80 ± 1.72	30.10 ± 9.54	54.10 ± 18.25	failure
10. Full Noise	N _c ≠ 0, N _s = 0	N _c ≠ 0, N _s ≠ 0	0.96	0.30 ± 0.64	27.50 ± 3.17	2.10 ± 1.51	4.00 ± 3.71	31.70 ± 8.15	69.10 ± 29.90	failure
11. Standardization	No	Yes	0.50	78.70 ± 3.93	62.70 ± 7.40	67.70 ± 8.99	79.60 ± 6.84	failure	41.80 ± 17.97	failure
12. Samples	n=1000	n = 20	0.92	failure	76.10 ± 4.72	failure	failure	failure	failure	error
13. Fixed support	No	Yes	0.87	failure	73.50 ± 4.18	failure	failure	failure	56.10 ± 8.25	failure

5% sparsity

Nodes d, samples n	SparseRC	NOTEARS	GOLEM
d = 200, n = 500	0	171	76
d = 500, n = 1000	0	377	114
d = 1000, n = 5000	0	639	44
d = 2000, n = 10000	0	171	time-out
d = 3000, n = 10000	0	1904	time-out

Non-public, real world gene

interaction data with interventions

CausalBench Challenge 2023 - Winners

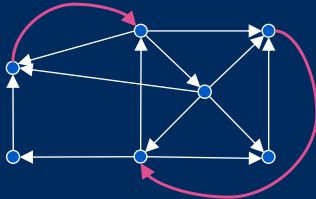
GRP 6000
1 Achilles Nazaret and Justin Hong
Columbia University

GRP 3000
2 Kaiwen Deng and Yuangfang Guan
University of Michigan

GRP 1500
3 Panagiotis Misiakos, Chris Wendler, Markus Püschel
ETH Zurich

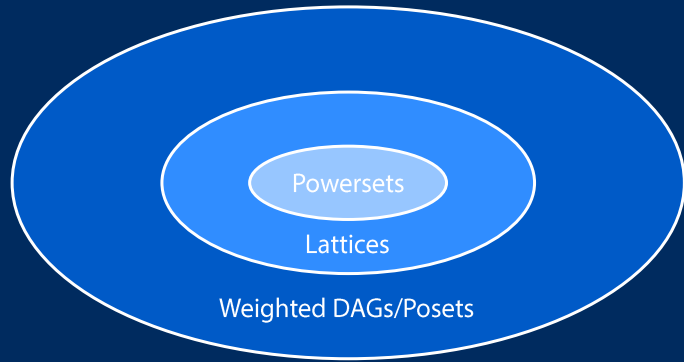
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Fourier basis for digraphs

Generalized the idea of cyclic boundary condition to arbitrary digraphs



Novel form of SP on DAGs and Posets

*shift-invariant
multiple shifts
interpretable
relates to causal reasoning*

Driven by insights from Algebraic Signal Processing

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